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ON THE SPECTRUM OF A GRAPH

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ABSTRACT

In this paper we discuss about the spectrum of a graph. We obtain the relations between a regular graph and its spectrum, and a complete graph and its spectrum, respectively. We obtain a bound for eigenvalues of an oriented graph with loops as a generalization of a non-oriented graph without loops. We prove that the maximum eigenvalue of a graph equals to its upper bound and its lower bound if and only if the graph is a complete graph and a regular graph, respectively.

1. The spectrum of regular graphs

Let G be a graph whose vertex-set VG is the set $\{v_1, v_2, \dots, v_n\}$ and whose edge-set EG is the subset of the set of unordered pairs of elements of VG . We call a graph with n vertices and m edges is an (n, m) non-oriented graph. A vertex-subgraph of G is a graph constructed by taking a subset U of VG together with all edges of G which are incident in G only with vertices belonging to U .

The adjacency matrix $A(G)$ of an (n, m) non-oriented graph G is an $n \times n$ symmetric matrix whose entries a_{ij} are given by

$$a_{ij} = \begin{cases} 1 & \text{if } \{v_i, v_j\} \in EG \\ 0 & \text{if } \{v_i, v_j\} \notin EG \end{cases} \quad (1.1)$$

The spectrum of an (n, m) graph G , $\text{Spec } G$, is the set of eigenvalues of $A(G)$ together with their multiplicities. Namely, if the distinct eigenvalues of $A(G)$ are $\lambda_1 > \lambda_2 > \dots > \lambda_s$, and their multiplicities are $m(\lambda_1), m(\lambda_2), \dots, m(\lambda_s)$, then we write the spectrum of a graph G by

$$\text{Spec } G = \left(\begin{array}{cccc} \lambda_1 & \lambda_2 & \cdots & \lambda_s \\ m(\lambda_1) & m(\lambda_2) & \cdots & m(\lambda_s) \end{array} \right). \quad (1.2)$$

We also write the maximum and minimum eigenvalues of $A(G)$ by $\lambda_{\max}A(G)$ and $\lambda_{\min}A(G)$, respectively. We use the notation $\lambda_{\max}(G)$ and $\lambda_{\min}(G)$ in place of $\lambda_{\max}A(G)$ and $\lambda_{\min}A(G)$, respectively.

Now let us consider the spectrum of a regular graph. A graph is said to be regular of degree k if each of its vertices has degree k . It is known that a regular graph G of degree k has $\lambda_{\max}(G)=k$, and $m(\lambda_{\max}(G))=1$ if G is connected.

LEMMA 1. An (n, m) graph G which has p connected components is regular of degree $(2m/n)$ if and only if $\lambda_{\max}(G)=(2m/n)$ and $m(\lambda_{\max}(G))=p$.

Proof. (\Rightarrow) By a suitable labelling of the vertices of G , the adjacency matrix $A(G)$ can be written in the partitioned form

$$A = \begin{bmatrix} A_1 & & 0 \\ & A_2 & \\ & & \ddots \\ 0 & & & A_p \end{bmatrix} \quad (1.3)$$

where submatrices $A_i(i=1, \dots, p)$ are corresponding to the adjacency matrices of a connected component G_i of G .

As G is a regular connected graph of degree $2m/n$ and $\lambda_{\max}(G_i)=2m/n$ and its multiplicity is 1, i.e. $m(\lambda_{\max}(G_i))=1$. As the eigenvalues of G consist of all eigenvalues of G_1, G_2, \dots, G_p , $\lambda_{\max}(G)=2m/n$ and its multiplicity is p .

(\Leftarrow) For any real $n \times n$ symmetric matrix X and for any real non-zero column n -vector \mathbf{z} , we call $\{(\mathbf{z}, X\mathbf{z})/(\mathbf{z}, \mathbf{z})\}$ be Rayleigh quotient and denote it by $R(X; \mathbf{z})$. Here (\mathbf{x}, \mathbf{y}) is the inner product of vector \mathbf{x} and \mathbf{y} . It is known that

$$\lambda_{\max}(X) \geq R(X; \mathbf{z}) \geq \lambda_{\min}(X) \quad \text{for } \forall \mathbf{z} \neq 0 \quad (1.4)$$

and the equality $R(X; \mathbf{z})=\lambda_{\max}(X)$ holds if and only if \mathbf{z} is an eigenvector corresponding to the eigenvalue $\lambda_{\max}(X)$.

Now let us put $\mathbf{z} = \overbrace{[11 \cdots 1]}'^n$, then we have

$$\lambda_{\max}(G) \geq R(A(G); \mathbf{z}) = \frac{2m}{n} \quad (1.5)$$

On the other hand we have $\lambda_{\max}(G)=2m/n$ by the hypothesis. Hence \mathbf{z} is an eigenvector corresponding to the eigenvalue $2m/n$, that is to say, $A\mathbf{z}=\{(2m/n)\}\mathbf{z}$. This implies each row sum of A is $2m/n$ and so G is a regular graph of degree $2m/n$.

Let a graph G has k components, than all k components of G are regular connected graphs of degree $(2m/n)$. Each component of G has $\lambda_{\max}(G)=2m/n$ whose multiplicity is one. Hence we have $k=p$.

2. The spectrum of the complete graph

The complete graph K_n has n vertices and each distinct pair is adjacent. It is known the spectrum of the complete graph K_n is

$$\text{Spec } K_n = \begin{pmatrix} n-1 & -1 \\ 1 & n-1 \end{pmatrix}. \quad (2.1)$$

LEMMA 2. If the spectrum of G is

$$\text{Spec } G = \begin{pmatrix} n-1 & -1 \\ 1 & n-1 \end{pmatrix}, \quad (2.2)$$

then G must be the complete graph K_n .

Proof. Let G be an (n, m) graph and the eigenvalues of $A(G)$ be $\lambda_1, \lambda_2, \dots, \lambda_n$ ($\lambda_1 \geq \lambda_2 \geq \dots \geq \lambda_n$). Then

$$\sum_{i=1}^n \lambda_i^2 = \text{tr}(A^2) = 2m \quad (2.3)$$

and

$$\sum_{i=1}^n \lambda_i = (n-1)^2 + (-1)^2(n-1) \quad (2.4)$$

$$\therefore m = \frac{n(n-1)}{2} \quad (2.5)$$

Hence it follows that G must be the complete graph K_n .

3. The lower and upper bounds for the maximum and minimum eigenvalues for a graph

LEMMA 3. For any (n, m) graph G with $n \geq 2$ and $m \geq 1$, we have

$$\lambda_{\max}(G) \geq 1, \quad -1 \geq \lambda_{\min}(G) \quad (3.1)$$

Proof. Any (n, m) graph G with $n \geq 2$ and $m \geq 1$ has at least one $(2, 1)$ vertex-subgraph G_1 . Let A_1 be the adjacency matrix of the vertex-subgraph G_1 , then the adjacency matrix A of G can be written in the partitioned form

$$A = \begin{bmatrix} A_1 & * \\ * & * \end{bmatrix}, \quad \text{where } A_1 = \begin{bmatrix} 0 & 1 \\ 1 & 0 \end{bmatrix} \quad (3.2)$$

Let \mathbf{x}_1 be a 2-vector which satisfies the condition $A_1\mathbf{x}_1=\lambda_{\max}(G_1)\mathbf{x}_1$ and \mathbf{z}_1 be a 2-vector which satisfies the condition $A_1\mathbf{z}_1=\lambda_{\min}(G_1)\mathbf{z}_1$. Let us put $\mathbf{x}=[\mathbf{x}_1, \overbrace{0 \cdots 0}^{n-2}]$ and $\mathbf{z}=[\mathbf{z}_1, \overbrace{0 \cdots 0}^{n-2}]$. Then

$$\lambda_{\max}(G_1)=R(A_1; \mathbf{x}_1)=R(A; \mathbf{x})\leq\lambda_{\max}(G) \tag{3.3}$$

$$\lambda_{\min}(G_1)=R(A_1; \mathbf{z}_1)=R(A; \mathbf{z})\geq\lambda_{\min}(G) \tag{3.4}$$

As

$$\lambda_{\max}(G_1)=1 \quad \lambda_{\min}(G_1)=-1, \tag{3.5}$$

we have

$$\lambda_{\max}(G)\geq 1 \quad -1\geq\lambda_{\min}(G). \tag{3.6}$$

LEMMA 4. If a connected (n, m) graph G with $n\geq 3$ is not the complete graph K_n , then

$$\lambda_{\max}(G)\geq 2^{1/2}, \quad -2^{1/2}\geq\lambda_{\min}(G) \tag{3.7}$$

Proof. If a connected (n, m) graph G with $n\geq 3$ is not the complete graph K_n , G has at least one $(3, 2)$ vertex-subgraph G_1 . A can be partitioned as follows:

$$A=\left[\begin{array}{c|c} A_1 & * \\ \hline * & * \end{array}\right], \quad \text{where } A_1=\begin{bmatrix} 0 & 1 & 1 \\ 1 & 0 & 0 \\ 1 & 0 & 0 \end{bmatrix} \tag{3.8}$$

As

$$\lambda_{\max}(G)\geq\lambda_{\max}(G_1) \tag{3.9}$$

$$\lambda_{\min}(G)\leq\lambda_{\min}(G_1) \tag{3.10}$$

and

$$\lambda_{\max}(G_1)=2^{1/2} \quad \lambda_{\min}(G_1)=-2^{1/2}. \tag{3.11}$$

Then we have

$$\lambda_{\max}(G)\geq 2^{1/2} \quad \lambda_{\min}(G)\leq -2^{1/2}. \tag{3.12}$$

4. A bound for the eigenvalues of an oriented graph

In this section, we consider an (n, m) oriented graph G with loops whose eigenvalues of the adjacency matrix A are all real numbers. The difference between non-oriented graph and an oriented graph is only its adjacency matrix A is symmetric or not.

THEOREM 1. Let an (n, m) oriented graph G has t loops and c cycles whose length are 2, then

$$\frac{t - \sqrt{(n-1)(2cn + nt - t^2)}}{n} \leq \lambda_i \leq \frac{t + \sqrt{(n-1)(2cn + nt - t^2)}}{n} \quad (4.1)$$

$(i=1, 2, \dots, n)$

Proof. By the hypotheses we have

$$\sum_{i=1}^n \lambda_i = \text{tr}(A) = t \quad (4.2)$$

$$\sum_{i=1}^n \lambda_i^2 = \text{tr}(A^2) = 2c + t \quad (4.3)$$

Let us put $x = [\lambda_2, \lambda_3, \dots, \lambda_n]'$ and $y = [\overbrace{1, 1, \dots, 1}^{n-1}]'$. Then

$$|(\mathbf{x}, \mathbf{y})| = \left| \sum_{i=2}^n \lambda_i \right| = |t - \lambda_1| \quad (4.4)$$

$$\|\mathbf{x}\| = \sqrt{\sum_{i=2}^n \lambda_i^2} = \sqrt{(2c + t) - \lambda_1^2} \quad (4.5)$$

$$\|\mathbf{y}\| = \sqrt{\sum_{i=2}^n 1^2} = \sqrt{n-1} \quad (4.6)$$

For these \mathbf{x} and \mathbf{y} , applying the Cauchy-Schwarz inequality we have

$$|t - \lambda_1| \leq \sqrt{(2c + t) - \lambda_1^2} \sqrt{n-1} \quad (4.7)$$

$$\therefore \frac{t - \sqrt{(n-1)(2nc + nt - t^2)}}{n} \leq \lambda_1 \leq \frac{t + \sqrt{(n-1)(2nc + nt - t^2)}}{n} \quad (4.8)$$

In a similar fashion, we can show that the inequality (4.8) holds for any $\lambda_i (i=2, 3, \dots, n)$.

This theorem can be reduced to a non-oriented graph with loops. The number of cycles whose length are 2 in an oriented graph equals to the number of edges in a non-oriented graph, i.e. $c=m$. By putting $c=m$ and $t=0$ in (4.8), we have

$$-\sqrt{\frac{2m(n-1)}{n}} \leq \lambda_i \leq \sqrt{\frac{2m(n-1)}{n}} \quad (i=1, 2, \dots, n) \quad (4.9)$$

Especially

$$\lambda_{\max}(G) \leq \sqrt{\frac{2m(n-1)}{n}}, \quad (4.10)$$

and this result coincides with the result which has already established.

5. The upper and lower bounds for the maximum eigenvalue of a graph.

Now let us consider a non-oriented graph again.

THEOREM 2. In an (n, m) graph G with $n \geq 1$ and without loops, we have

$$\frac{2m}{n} \underset{(1)}{\leq} \lambda_{\max}(G) \underset{(2)}{\leq} \sqrt{\frac{2m(n-1)}{n}}, \quad (5.1)$$

where the equality (1) holds if and only if G is a regular graph of degree $(2m/n)$ and the equality (2) holds if and only if G is the complete graph K_n .

Proof. The equality of (1) is clear from Lemma 1.

We will prove that the equality of the Cauchy-Schwarz inequality

$$|(\mathbf{x}, \mathbf{y})| \leq \|\mathbf{x}\| \cdot \|\mathbf{y}\| \quad (5.2)$$

holds if and only if $\lambda_{\max}(G) = \sqrt{2m(n-1)/n}$ is satisfied.

Let us put $\mathbf{x} = [\lambda_2, \lambda_3, \dots, \lambda_n]'$ and $\mathbf{y} = [1, 1, \dots, 1]'$, then the equality of (5.2) holds if and only if $\mathbf{y} = a\mathbf{x}$ is satisfied, that is when $\lambda_2 = \lambda_3 = \dots = \lambda_n (\equiv \lambda)$ is satisfied. As the graph G has no loops,

$$0 = \text{tr}A = \lambda_1 + (n-1)\lambda = \lambda_{\max}(G) + (n-1)\lambda \quad (5.3)$$

$$\therefore \lambda_{\max}(G) = -(n-1)\lambda, \quad (5.4)$$

we have

$$|(\mathbf{x}, \mathbf{y})| = \left| \sum_{i=2}^n \lambda_i \right| = |(n-1)\lambda| = |-\lambda_{\max}(G)|. \quad (5.5)$$

As the graph G has m edges,

$$2m = \text{tr}(A^2) = \sum_{i=1}^n \lambda_i^2 = \lambda_1^2 + (n-1)\lambda^2 = \lambda_{\max}^2(G) + (n-1)\lambda^2 \quad (5.6)$$

$$\therefore (n-1)\lambda^2 = 2m - \lambda_{\max}^2(G), \quad (5.7)$$

we have

$$\|\mathbf{x}\| \cdot \|\mathbf{y}\| = \sqrt{(n-1)\lambda^2} \cdot \sqrt{n-1} = \sqrt{2m - \lambda_{\max}^2(G)} \sqrt{n-1}. \quad (5.8)$$

From (5.5) and (5.8),

$$|-\lambda_{\max}(G)| = \sqrt{2m - \lambda_{\max}^2(G)} \sqrt{n-1} \quad (5.9)$$

$$\therefore \lambda_{\max}(G) = \sqrt{\frac{2(n-1)m}{n}} \quad (5.10)$$

$$\therefore \lambda = -\sqrt{\frac{2m}{(n-1)n}} \quad (5.11)$$

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Hence the spectrum of G must be

$$\text{Spec } G = \begin{pmatrix} \sqrt{\frac{2m(n-1)}{n}} & -\sqrt{\frac{2m}{n(n-1)}} \\ 1 & n-1 \end{pmatrix}. \quad (5.12)$$

As $m \leq n(n-1)/2$, it must

$$-\sqrt{\frac{2m}{n(n-1)}} \geq -1. \quad (5.13)$$

On the other hand, by Lemma 3,

$$-1 \geq \lambda_{\min}(G) = -\sqrt{\frac{2m}{n(n-1)}}. \quad (5.14)$$

It must

$$-\sqrt{\frac{2m}{n(n-1)}} = -1 \quad (5.15)$$

$$\therefore m = \frac{n(n-1)}{2}, \quad (5.16)$$

and the graph G must be the complete graph K_n . Hence we can conclude that the equality of (2) holds if and only if G is the complete graph K_n .

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