Title	An interactive debugging system composed of a minicomputer and a microprocessor
Sub Title	
Author	Okada, Kenichi(Matsuo, Taiki) 松尾, 泰樹(Kitagawa, Misao) 北川, 節
Publisher	慶応義塾大学工学部
Publication year	1980
Jtitle	Keio engineering reports Vol.33, No.10 (1980. 12) ,p.131- 150
JaLC DOI	
Abstract	This paper describes a debugging system which has been developed on such a combined system of a minicomputer and a microprocessor that is microprogrammable for the user, and aims at an effective debugging of the errors that will be detected during the execution of the program written in a low-level language. The multiprocessor organization, adoption of firmware monitor and special hardwares yield such advantageous features as bilateral tracing, procedure extraction, eight kinds of event monitors, etc., which come to be effective for debugging. Debugging is performed in art interactive mode so that the program can be tested and modified easily by various kinds of editing and debugging commands. This system has four modes, that is, assemble, edit, debug, and test, which allow it to perform the debug processing consistently.
Notes	
Genre	Departmental Bulletin Paper
URL	https://koara.lib.keio.ac.jp/xoonips/modules/xoonips/detail.php?koara_id=KO50001004-00330010- 0131

慶應義塾大学学術情報リポジトリ(KOARA)に掲載されているコンテンツの著作権は、それぞれの著作者、学会または出版社/発行者に帰属し、その権利は著作権法によって 保護されています。引用にあたっては、著作権法を遵守してご利用ください。

The copyrights of content available on the KeiO Associated Repository of Academic resources (KOARA) belong to the respective authors, academic societies, or publishers/issuers, and these rights are protected by the Japanese Copyright Act. When quoting the content, please follow the Japanese copyright act.

AN INTERACTIVE DEBUGGING SYSTEM COMPOSED OF A MINICOMPUTER AND A MICROPROCESSOR

KEN'ICHI OKADA, TAIKI MATSUO and MISAO KITAGAWA

Dept. of Instrumentation Engineering, Keio University, Yokohama 223, Japan

(Received July 15, 1980)

ABSTRACT

This paper describes a debugging system which has been developed on such a combined system of a minicomputer and a microprocessor that is microprogrammable for the user, and aims at an effective debugging of the errors that will be detected during the execution of the program written in a low-level language. The multiprocessor organization, adoption of firmware monitor and special hardwares yield such advantageous features as bilateral tracing, procedure extraction, eight kinds of event monitors, etc., which come to be effective for debugging. Debugging is performed in an interactive mode so that the program can be tested and modified easily by various kinds of editing and debugging commands. This system has four modes, that is, assemble, edit, debug, and test, which allow it to perform the debug processing consistently.

1. Introduction

In contrast to the drastic reduction of the recent hardware prices, software cost in the computer system tends to rather increase its rate, since the software productivity will not largely be improved. [1] One of the factors that suppress the improvement of software productivity lies in the difficulty of program debugging; [9] especially, debugging the errors detected in the execution of programs is an urgent problem, and thus many researches for the debuggers have been studied up to the present. [2] [3] [4] [5] [6] [8].

It can be considered that the requirements for such a debugger include (1) an interactive processing, (2) time-sharing processing, (3) high-performance efficiency, (4) acquisition of high-level information, (5) feasible usage, and so forth. At the time of developing the currently proposed debugging system that can satisfy the above requirements, the following principles are established:

- (1) Its object is laid in the assembly language, and is to debug the errors detected in the execution of user program.
- (2) It is realized on the multiprocessors to distribute their functions and also to suppress the occurrence of secondary errors which might be made through debugging.
- (3) It should employ a firmware monitor.
- (4) The debugging is proceeded in an interactive mode, and thus the timesharing system is not taken into account because the object is limited to a minicomputer only.

Following the above statement, a debugging system which is realized on the combined system composed of a HITAC-10 minicomputer and a mCOM-16 microprocessor is described, which is shown in Figure 1, where the console for debugging use and various I/O units are also illustrated. [7]



Figure 1. Hardware configuration.

2. Debugger

The debugger's software is composed of the processing programs of both the minicomputer side (nearly 4500 steps of the HITAC-10 assembly language) and the microprocessor side (about 2500 steps of the mCOM-16 microprogram), and collects and displays the debugging information obtained in the execution of the program in accordance with the commands delivered from the user. The user program is stored in the main memory of the mCOM-16 and carried out through the firmware monitor (three kinds of firmware monitors are registered in the disk and transferred to the mCOM-16 control memory). Since the mCOM-16 is controlled and managed by the HITAC-10, the user needs not be conscious of the existence of the mCOM-16. On the other hand, the processing program in the HITAC-10 is

divided into three; namely, (1) the command processing part which processes and analyzes the user's commands, (2) the debugging information part which arranges and displays the information obtained from the firmware monitor, and (3) the I/O service part which is prepared for the mCOM-16 providing no I/O units.

2.1 Firmware Monitor

For the purpose of collecting effective and high-level information without introducing any disturbances at the debug processing, the following are set as the monitor:

- (1) Execution of instructions by the specified number of steps.
- (2) Execution of the specified instructions by the specified number of times.
- (3) Execution of the instruction at the specified location by the specified number of times.
- (4) Execution of instructions at the specified area of the memory.
- (5) Read out from the specified location or area.
- (6) Write into the specified location or area.
- (7) Comparison of the contents of the memory with the constants.
- (8) Comparison of the contents between two different memory locations.

The multiple monitors can be designated simultaneously for those events, and even when only one of them is specified during the execution of user program, the execution is halted to display it to the user. Thus the user can recognize the executing state at the time when the individual conditions have been satisfied.

Furthermore, the following are prepared as the back-up functions in order to proceed the debug processing smoothly:

(1) Restoring function of the initial state:

Since the user program stored in the disk is transferred to the main memory of the mCOM-16, and at the same time the initialization of the firmware monitor is carried out, the reexecution of the user program can be made from the initial state arbitrarily.

(2) Read and write functions for the registers and memory:

The contents of all the registers and memory of the emulated machine can be read or written.

(3) Procedure extraction function:

A series of processes which affect a particular location are taken out as a unit of procedure. Namely, this is such a function that since a sequence of instructions composed of a program has a pattern that it may start with a transfer instruction and then end at another transfer instruction after processed by some arithmetic instruction(s), such a sequence is picked up from the program to display as a unit of procedure. It will be a powerful tool for the user in a debugging operation to recognize a series of instructions that affect the location, the value of which is found to be inadequate.

(4) Backtrack function:

Backtracking, which proceeds the execution of user program in the reverse direction, can be realized by stacking and monitoring the data information, which comprises the contents of memory and registers that

will be destroyed during the forward execution of the program and the address information concerning the program flow at the Jump, Call, Return instructions and the like in the ring stack. The number of steps that can be backtracked amounts to approximately 1000 steps at present, which will be determined by the ring stack capacity.

(5) Trace function:

This function provides a means of displaying the contents of the respective registers and effective addresses everytime when the specified operations, specified operands, or the instructions in the specified area are carried out, and consequently can catch the consecutive state transitions of user program. During the trace, both collection of the trace information due to the firmware monitor on the mCOM-16 and its display to the user through the HITAC-10 are performed in parallel.

(6) I/O change function:

This permits the change of the I/O units used in the user's program without altering its construction. Presently the TTY and PTR are provided.

(7) Disk dump function:

This allows the user program to be saved in the disk at any time and to return to that point, and comes to be effective to realize a check point and restart function.

These functions as well as the designation of events to be monitored are carried out using the commands, and the HITAC-10 generates the Reference Table which will be mentioned later in accordance with the given commands. The firmware monitor refers to the Reference Table and proceeds the execution of user program, collecting the useful information only. Events are checked at the following three processing parts:

- (1) Fetch part, where the execution address, the number of steps, and the kind of instructions are checked.
- (2) Address computation part, where the memory access is checked.
- (3) Execution part, where the executed results are checked.

The result of individual checking, if an event specified by the user comes into existence, will be informed to the HITAC-10, which then displays the generated event, the contents of individual registers, and the contents of consecutively seven locations of the program placing the approved location in the center of them.

The following three kinds of runs are registered in the firmware monitor: (1) Emulation run, which proceeds the execution of user program while monitoring the event specified by the user's commands; (2) preparation run, where the stacking operation for backtrack use is added to the emulation run; and (3) back run, which executes the program in the reverse direction while monitoring the event. Those are stored in the disk and transferred to the control memory of the mCOM-16 through user commands, and also the mapping memory is rewritten. Figures 2 and 3 depict these three kinds of monitor flows.

2.2 Commands and Reference Table

For the purpose of easy input and of getting abundant, useful information, the commands listed in Table 1 are prepared. As the operand of command, one

can freely use the symbolic names in the user program, the direct address designation by decimal or hexadecimal representation, and the mnemonic names of registers. Multiple commands input from the console for debugging use are rearranged and compiled to generate the Reference Table.



Figure 2. Flow of the firmware monitor (ER, PR).



Figure 3. Flow of the firmware monitor (BR).

Initialize Dump to disk Dump to memory Set register Set memory Get register Get memory Break point Program counter
Dump to disk Dump to memory Set register Set memory Get register Get memory Break point Program counter
Dump to memory Set register Set memory Get register Get memory Break point Program counter
Set register Set memory Get register Get memory Break point Program counter
Set memory Get register Get memory Break point Program counter
Get register Get memory Break point Program counter
Get memory Break point Program counter
Break point Program counter
Program counter
Step
Operation code
Operand refer
Operand modify
Compare with constant
Compare with variable
Extract
I/O change
Cancel I/O change
Reverse assemble
Trace
Untrace
Emulation run
Preparation run
Back run
Finish
Go to command waiting
Cancel current command

Table 1. Command of the debugger.

r:	range
	address—address
a:	address
	decimal number, hexadecimal number, or label
n:	register name
	PC, AC, EC, CAR, or IR
v :	value
	decimal number, or hexadecimal number
o :	operator
	EQ, GT, GE, LT, LE, or NE
c:	operation code
i :	I/O device name
	PTR or TTY
-:	repeatable

.



Figure 4. Format of the Reference Table.

The Reference Table determines the operation of the firmware monitor and has 10 kinds of formats which keep the following information, as given in Figure 4:

- (1) Number of steps to be executed.
- (2) Location or area where the execution is prohibited.
- (3) Location or area to be traced.
- (4) Number of times of executing a location till the time of activating the breakpoint and/or break mechanism.
- (5) Specified number of accesses to a location and the necessity of tracing.
- (6) Memory area where the access is to be monitored.
- (7) 'Compare' operator and two locations to be compared or a location and a constant.
- (8) Specified number of times of executing macroinstructions and the necessity of tracing.
- (9) A location to be noticed by the procedure extraction function.
- (10) Names of two I/O devices to be changed each other.

In the above stated Reference, (1), (9), and (10) include only one item each; (8) will have the same number of items as that of the macroinstructions; and (2)

through (7) will keep the necessary number of entries for themselves. This Reference Table will be transferred to the special area of the main memory of the mCOM-16 through the execution (run) commands of the user program.

2.3 Processing of the Minicomputer Side

While the firmware monitor on the mCOM-16 executes the user's program, the processing that the minicomputer performs in the debugger includes the following:

- (1) Input of the user commands from the console.
- (2) Analysis of the commands, generation and updating of the Reference Table.
- (3) Read out from and write into the registers and the memory of the mCOM-16.
- (4) Transfer of the microprogram to the control memory of the mCOM-16 and its initiation.
- (5) I/O services for the mCOM-16.
- (6) Rearrangement and display of the debugging information resulted from the firmware monitoring.
- (7) Mode change of the system.

Figures 5 and 6 show the flow of processing by the minicomputer of the



Figure 5. Flow of the debugger in debug mode.



Figure 6. Flow of the debugger in test mode.

debugger.

This system has the four operation modes as follows:

- (1) Assemble mode that converts the user's program to the object one.
- (2) Edit mode that performs the program compilation and the file processing.
- (3) Debug mode that arranges the condition for the debug processing and displays the debug information.
- (4) Test mode that the firmware monitor performs the execution of user program.

Table 2 lists up the contents of the respective memory area in those four modes. On the HITAC-10, the assembler for debugging use runs in assemble mode, the text editor runs in edit mode, and the debugger runs in either debug or test mode; the distinction between them depends upon whether or not the firmware monitor in the control memory of the mCOM-16 is operating.

The role of the assembler for debugging use is to convert the user program to the object one and it is provided as the debugging version after having modified the macro-assembler delivered from the manufacturer. The principal modification is in that the address designation by symbols while debugging the program can be performed by way of saving the symbol table in a safety area at the end of assembling.

	Edit Mode	Assemble Mode	Debug Mode	Test Mode
HITAC-10 Main Memory	Editor	Assembler	Debugger	Debugger
Common Area	_	_	Reference Table	Reference Table
mCOM-16 Main Memory	_	_	Object	Object
mCOM-16 Control Memory	_	. —	-	Firmware Monitor

Table 2.	Contents of	the	specified	memories	for	four	system	modes.



Figure 7. Mode transition.

The text editor has the function of various kinds of compilation and file processing, and is capable of linking the assembler with the debugger.

Figure 7 gives the operation and transition of these four modes, where the transition from assemble mode to edit is initiated automatically after terminating the assembling; the transition from test mode to debug is due to the generation of a specified event, and the other transitions are generated by the user's commands. In debug and test modes the following data transfer is carried out between the HITAC-10 and the mCOM-16.

- (1) From HITAC-10 to mCOM-16,
 - (a) object codes of the user's program (in debug mode),

- (b) firmware monitor (in debug mode),
- (c) data to the mapping memory (in debug mode),
- (d) Reference Table (in debug mode), and
- (e) I/O service information (in test mode).
- (2) From mCOM-16 to HITAC-10,
 - (a) debug information (in test mode), and
 - (b) I/O service request (in test mode).

These data transfer will be performed through the two buffer registers connecting the I/O Bus of the HITAC-10 with the U-Bus of the mCOM-16.

3. Examples of Debug Processing

This chapter describes some examples concerning the actual debug processing by making use of the results obtained from the execution of a sample program (see Appendix) through the debugger. The program is to convert a 4-digit hexadecimal (4 bytes) number of the character format to the decimal, using the TTY as an I/O unit. Each input hexadecimal digit of 1 byte is converted to a 4-bit hexadecimal digit after echobacking and stored in a word of 16 bits packed to the lefthand side, where the most significant bit is considered to be the sign. The 4bit hexadecimal number thus obtained is output after converting it to the character type decimal number. The following lists are the output results obtained from the execution of the sample program with several kinds of commands picked out of those in this debugger. (\$\$ indicates the mode change, ** gives the system message in debug mode, mark # shows the address where the program is stopped after the HLT instruction has been carried out or some condition has been satisfied.)

(1) From edit to debug mode by Z command. First, initialization is carried out. (See List 1.)

List 1. \$\$ TTY - EDITOR \$\$ >Z \$\$ HITAC-10 DEBUGGER \$\$ #I

(2) The output results obtained from the normal execution without any input of the commands for event monitor use. (See List 2.)

The Program Counter is set to the starting location and the user program is executed by the ER command. Since the input routines from the TTY (X'2005' to X'200A') is in the user program, the operation will be at the state of queue. When 0800 is input at this point, =_02048 (_______ indicates a space) is output by the user program, and the debugger informs the user that the HLT instruction has been carried out. The indicated information includes the following: the location and the instruction in it at the time its execution ended, the contents of Accumulator (AC), Extended Accumulator (EC), Carry Register (CAR), Effective Address (EA) and its contents ((EA)) as well as the location where the instruction just finished

List 2.

#SR PC X"2000" #ER 0800 = 02048** I'VE JUST EXECUTED "HLT" INSTRUCTION. ** LOC IR AC EC CAR EA (EA) 2033 6F00 000A 0000 0 KCT 2030 5266 ADR2 *+054 #2066 2031 5259 КСТ *+040 #2059 2032 422E B *-004 #202E 6F00 HLT * 2033 0A6B 2034 X20 *+055 #206B 1 2035 385E ST BUFF *+041 #205E 2036 4200 В BGN *-054 #2000 *** CONTINUE ? ** List 3. #CC M7+1 EQ 0 ** THE SPECIFIED LOCATION HAS THE APPOINTED VALUE NOW.** #BP BGN+1 #FR ** I'VE'JUST DONE THE SPECIFIED LOCATION.** ΕC CAR EA LOC IR AC (EA) 2001 5259 205E 0000 0 2059 FFFB => FFFA 472 #01C4 1FFE F1C4 STE Ζ 1FFF 9156 SRA 342 #0176 BGN 2000 4A37 BAL INIT *+055 #2037 * 2001 5259 KCT *+088 #2059 KÇT KCT *+087 #2059 *+086 #2059 2002 5259 5259 2003 2004 KCT ADR2 *+098 #2066 5266 #CC M7+1 EQ 0 #ER FFFB ** THE SPECIFIED LOCATION HAS THE APPOINTED VALUE NOW .** LOC IR AC ЕC CAR EA (EA) 2059 FFFF => 0000 201D 5259 FFFB 0000 0 SLL 004 #0004 201A 8804 201B 325B 0 HEXA *+064 #205B ST 2010 3A5C PACK *+064 #205C * 201D 5259 KCT *+060 #2059 READ *-025 #2005 201E 4205 В 201F 6864 КРА 2020 4A51 BAL NEGT *+049 #2051 #GR AC CONTENT OF AC ;#FFFB

#GM PACK CONTENT OF LOCATION #205C IS # FFFB

List 4. #BP X"2021" **#T NEGT, NEGT+6** #ER LOC IR AC EC CAR EA (EA) 2052 2050 0004 0000 0 2050 FFFF *-002 #2050 2052 2050 X EC CAR EA (EA) LOC IR AC 2053 3A5E 0004 0000 0 205C FFFB => 0004 2053 3A5C ST PACK *+009 #205C AC EC LOC IR CAR EA (EA) 2050 0004 => 0005 2054 525C 0004 0000 0 2054 5250 KCT PACK *+008 #205C ЕĈ LOC IR AC CAR EA (EA) 206A 002D 2055 0A6A 002D 0000 0 2055 006 L X2D *+021 #206A CAR EA AC EC (EA) LOC IR 205E 0020 => 002D 2056 3A5E 002D 0000 0 2056 3A5E ST BUFF *+008 #205E AC EC CAR EA (EA) LOC IR 2057 4651 002D 0000 0 2021 0A5A В NEGT *-006 #2051 Ī 2057 4651 ** I'VE JUST DONE THE SPECIFIED LOCATION.** EC CAR EA 100 IR AC (FA) 2021 0A5A 003D 0000 0 205A 003D 4205 READ *-025 #2005 201E B KPA 201F 6864 2020 4A51 BAL NEGT *+049 #2051 2021 0A5A 1 EQ *+057 #205A WRTE *+038 #2048 BAL 2022 4948 2023 4A3D CUHD *+026 #203D BAL

its execution and the inversely assembled results of seven locations (2030–2036). At the end of the HLT instruction, the 'CONTINUE' message is issued and will start the execution of program from the next location when 'C' is input from the TTY. In case where the characters except 'C' are input, the debugger will be at the state of accepting the commands of debug mode.

Ι

*+236 #2110

Ν

2024

2710

(3) List 3 is an example to confirm that the 4-byte input data are packed and stored in the PACK location.

Since location M7+1 is a counter of the input characters, the command to halt is input when the contents of location M7+1 has come to be zero

List 5.

#BP X"202D"					
#ER = ** I'VE JUS	T DONE THE	SPECIFI	ED LOCATIO	N xir xir	
LOC IR A 2020 4A37 0	C EC C 035 0000 0	CAR EA 2037 :	(EA) 2001 => 20	2E	
202A 202B 202C	000A 4A3D 0001	BAL	CUHD	*+018	#2030
* 202D 202E 202F 2030	4A37 ØE66 4A48 5266	BAL BAL KCT	INIT I ADR2 WRTE	*+010 *+056 *+025 *+054	#2037 #2066 #2048 #2066
#GM BUFF,BU CONTENT OF	FF+5 LOCATION	#205E IS	# 002D		*2000
CONTENT OF CONTENT OF	LOCATION LOCATION	#205F IS #2060 IS #2061 IS	# 0030 # 0030 # 0030		
CONTENT OF CONTENT OF	LOCATION	#2062 IS #2063 IS	# 0030 # 0035		

(input of 4 characters finished). The message that the conditions have already been satisfied is issued (which will be checked before executing the CC and CV commands), so that the program is executed to the location where the initialization of it completed. Here again the command to halt is input when the contents of location M7+1 come to be zero and then the program is executed. As the the result, such a message is output that the conditions are satisfied after the input of hexadecimal date 'FFFB' of 4-byte form. It can be confirmed that the input data in the character format are stored in the form of hexadecimal numbers after reading the contents of the AC and PACK locations.

(4) List 4 gives an example to confirm the state transition that the minus datum FFFB stored in the PACK location is converted to a form of 2's complement.

While the program is executed to location X'2021', the command to trace locations NEGT to NEGT+6 (conversion subroutine to the 2's complement) is input. It can be confirmed below that the contents 'FFFB' of the PACK location comes to be 0005 and the minus code X'2D' is embedded in the BUFF location.

- (4) List 5 is an example to confirm that the converted data to the decimal numbers are stored in locations BUFF to BUFF+5 in the character format. Such a command is input that executes the program till the time when embedding the character type decimal data to the BUFF is completed. It can be confirmed that the contents of locations BUFF to BUFF+5 are read and '-00005' is embedded.
- (6) Examples of the re-initialization setting, modify, step, and extract.

By making use of the I command, it is possible to execute the program again from its initial state at any arbitrary point of it. The program can

be executed by the input of the command to halt and by proceeding the stacking for the back trace use after the WORK location has been modified. At the state of input queue (due to the user program), when '8' is input, the message is output, whicn indicates that the read-in data X'38' has been stored in the WORK location. (See List 6.)

The following shows the results after executing 4 steps of the program

List 6.1. #I #SR PC X"2000" #OM WORK #PR 8 ** I'VE MODIFIED THE SPECIFIED LOCATION. ** LOC IR AC EC CAR EA (EA) 200E 3A67 0038 00B8 1 2067 0000 => 0038 200B 4048 BAL WRTE *+061 #2048 200C 8809 SLL 009 #0009 200D 8009 SRL 009 #0009 * 200E 3A67 ST WORK *+089 #2067 200F 1969 ς X40 *+090 #2069 2010 6864 KPA 2011 4215 R *+004 #2015 List 6.2. #ST 4 #PR ** I'VE DONE THE NUMBER OF STEPS YOU ORDERED. ** 1.00 IR AC EC CAR EA (EA) 2015 0A67 0038 00B8 1 2067 0038 2012 ØR67 L WORK *+085 #2067 A *+085 #2068 1268 2013 4216 В *+002 #2016 2014 2015 ·0A67 L WORK *+082 #2067 012 #0000 2016 880C SLL 2017 800C SRL 012 #000C ST 2018 3A5B HEXA *+067 #205B List 6.3. #E WORK #BR LOC IR AC EC CAR EA (EA) 2009 7074 205E 00B8 0 2009 С 7074 RTI 2000 8809 009 #0009 SLL SRL 009 #0009 200D 8009 WORK *+089 #2067 200E 3A67 ST

while stacking.

The following gives the results obtained after extracting a series of instructions that affect the WORK location while restoring the program inversely from the X'2015' location.

The extract instruction halts at the time when restoring the transfer instruction that affects the specified location. Accordingly, this example shows that the contents of the AC, EC, and CAR were X'205E', X'00B8', and 0, respectively, before executing the RTI C instruction.

To edit mode by the F command that follows,

List 6.4.

#F

\$\$ TTY - EDITOR \$\$

4. Conclusion

Since the program used in the previous examples of debug processing is so much simple and required some hand manipulation of the user that the reduction of execution efficiency by the firmware monitor has not been found at all. Therefore, when a program that has no I/O instructions to compute the factorial of 1000 is carried out under the following three conditions, the execution time is as follows:

- (1) Execution by the HITAC-10-42.0 msec.
- (2) Execution by the emulator-420.1 msec.
- (3) Execution by the firmware monitor—2032.0 msec.

These figures indicate that the result by the firmware monitor introduced approximately 50 times of speed reduction composed with that by the hardware, but the principal cause of it is due to the low performance of the microprocessor mCOM-16 and also the lack of the hardware such as multiplication, division, shifter, or the like. The evaluation of the overhead by means of supplementing the debugging function should be composed with the usual emulator, where the speed reduction amounts to about 5 times. In fact when this system was operated, the processing time for debugging information at the user side has come to be extremely large, and thus the decrease in execution speed of the program in the minicomputer does not give much trouble to the user. The most essential factor is that the profitable debugging information is obtained and thus the modification and testing of the program can easily be carried out. This system provides a simple mode change, many edit functions as well as monitor ones, which have all used to accomplish the purpose of debugging from the standpoint of improving the overall performance of the system.

REFERENCES

[1] ВОЕНМ, В. W., (1976): Software engineering, IEEE Trans. on Computers, Vol. C-25, No. 12, pp. 1226-1240.

- FLANAGAM, H., (1973): Program debugging system, IBM Technical Disclosure Bulletin, Vol. 16, No. 7, pp. 2322-2329.
- [3] GOLDBERG, J., COOPERBAND, A., and GALLENSON, L., (1978): PRIM System—A Framework for emulation-based debugging tools, National Computer Conference, pp. 373-377.
- [4] GRISHMAN, R., (1970): The debugging system AIDS, SJCC, pp. 59-64.
- [5] OKADA, K, and KITAGAWA, M., (1978): BRAKEMAN—A debugging hardware system, KEIO Engineering Reports, Vol. 31, No. 13, pp. 139-149.
- [6] OKADA, K., YOKOHAMA, T., and KITAGAWA, M., (1979): A dynamic debugging system by multiple processors (in Japanese), Technical Reports of Information Processing, IEE of Japan, IP-79-45, pp. 47-56.
- [7] OKADA, K. and KITAGAWA, M., (1980): A general-purpose experimental computer system characterized by its architecture changeability, KEIO Engineering Reports, Vol. 33, No. 5.
- [8] SAKAMURA, K., KITAFUSA, H., TAKEYARI, Y., and AISO, H., (1977): A debugging machine (in Japanese), Trans. IECE of Japan, Vol. J60-D, No. 9, pp. 671-678.
- [9] SCHWARZ, J.T., (1970): An overview of bugs, Debugging Techniques in Large Systems, PRENTICEHALL, INC., Englewood Cliffs, New Jersey.

Appendix

0800	00001 0800 <u>G</u>				
		1111. 1 1	CONVERT	HEXADECI	CIMAL TO DECIMAL /
2000 2000 2001 2002 2003	00001 2000G 4A37a 5259a 5259a 5259a	BGN	ORG BAL KCT KCT KCT KCT	X"2000' INIT M7+1 M7+1 M7+1 M7+1	9 U
2004 2005 2006 2007 2008	52660 6D000 70620 70610 42070	READ	KCT RIM STI KTI B	ADR2 *-1	READ HEXADECIMAL NUMBER
2009 200A 200B 200C 200C	7.0740 60000 48480 88090 80090 30670		SIM BAL SLL SRL	WRTE 9 9 NORK	 ECHO BACK REMOVE PARITY
200F 2010 2011 2012 2013 2014	19690 68640 42150 09670 12680 42160		S KPA B L A B	*+4 WORK WORK+1 *+2	✓ A B C D E F ?
2015 2016 2017 2018 2019 2019	0A670 880C0 800C0 3A5B0 0A5C0 88040	•	L SLL SRL ST L SLL	WORK 12 12 HEXA PACK 4	REMOVE ZONE CODE
201B 201C 201D 201E 201E	325B0 3A5C0 52590 42050 68640		O ST KCT B KPA	HEXA PACK M7+1 READ	STORE INTO ONE WORD
2020	4A510 2A540		BAL	NEGT FO	✓ WHEN NEGATIVE
2022 2023 2024 2025 2026 2026 2027	4A48a 4A3Da 2710a 4A3Da 03E8a 4A3Da		BAL BAL DC BAL DC BAL	WRTE CUHD 10000 CUHD 1000 CUHD	<pre> PRINT CHAR. = </pre>
2028 2029 202A 202B 202C 202C	00640 4A3D0 000A0 4A3D0 00010 4A370		DC BAL DC BAL DC BAL	100 CUHD 10 CUHD 1 INIT	
202E 202F 2030 2031 2032 2033	0E660 4A480 52660 52590 422E0 6E000		L ,I BAL KCT KCT B HIT	ADR2 WRTE ADR2 M7+1 *-4	PRINT DECIMAL NUMBER
2033 2034 2035 2036	0A6B0 3A5E0 42000		L ST B	X20 BUFF BGN	

2037 2038 2039 203A 203B 203C	00000 0A580 3A590 0A650 3A660 46370	INIT	DC L ST L ST B,I	M7 M7+1 ADR1 ADR2 *-5	/	INITIALIZE
203D 203E 203F 2040	00000 88000 C25C0 EE3D0	CUHD	DC SLL LE D ,I	PACK *-3	/	CONVERT
2041 2042	3A5CƏ A800Ə	÷	ST SLDL	PACK	1	REMAINDER> PACY
2043 2044 2045 2046 2047	125Də 3E66ə 5266ə 523Də 463Də		A ST ,I KCT KCT B ,I	X30 ADR2 ADR2 CUHD CUHD CUHD	,	QUOTIENT> (ADR2)
2048 2049	00000 60000	WRTE	DC		/	PRINT ONE CHAR.
204A	70860	2 ⁴⁴ 1	WTO			
2048 204C	42482		B	*-1		
204D 204E	70820 - 60000	• •	CTO SIM			
204F	46480	,	B ,I	WRTE		
2050	FFFFØ		DC	X"FFFF	1	
2051	2050a	NEGI	DC X	NEGT-1		
2053	34500		ST	PACK	_	And the second
2054 2055	525Ca 8666a		KCT	PACK X2D	/	TWO'S COMPLEMENT
2056	3A2E9		ST T	BUFF		
2007	46010	/	B *1	NEUI		
2058	FFF90 0001K	M7	DC	-7		
205A	003D0	EQ	DC	Ç"="		
205B 205C	0001K 00000	PACK	DC.	1		
205D	00300	X30	DC	X"30"		
205E 205F	00200 0005K	BUEE	DC DS	X"20" 5		
2064	00000		DC	-X"0A"		
2065 2066	205E0 0001K	HDR1 ADR2	DS	BUFF 1		
2067	0001K	WORK	DS	1		
2068 2069	00090 00400	X40	DC	9 X"40"		
206A	002D0 00200	X2D	DC	X"2D"	1	MINUS -
2000	F352I	A20	DC	A ZU	1	DEHUE
00B9 2289	0089G 3E000					
00BA	3FC00	- 41				
009F	7F80I 009FG					
009F	00000					
	00001 2000h		END	BGN		e e e e e e e e e e e e e e e e e e e

• • •

150